

TD 6: BDDs

Exercise 1 (Some BDDs). Draw the reduced BDDs for the following functions, using the order of your choice on the variables $\{x_1, x_2, x_3\}$:

1. $(x_1 \Leftrightarrow x_2) \vee (x_1 \Leftrightarrow x_3)$,
2. the majority function $m(x_1, x_2, x_3)$: its value is 1 iff the majority of the input bits are 1's,
3. the constant sum function $s_c(x_1, x_2, x_3)$ for $c = 1$: its value is 1 iff $c = \sum_{i=1}^3 x_i$,
4. the hidden weighted bit function $h(x_1, x_2, x_3)$: its value is that of variable x_s , where $s = \sum_{i=1}^3 x_i$ and x_0 is defined as 0.

Exercise 2 (Symmetric Functions). A *symmetric function* of n variables has the same value for all permutations of the same n tuple of arguments. Clearly, all variable orderings lead to the same reduced BDD size for symmetric functions.

Show that a reduced BDD for a symmetric function has at most $\binom{n+2}{2}$ nodes.

Exercise 3 (Counting Solutions). Write a (nearly) linear time algorithm for counting the number of solutions of a boolean function f represented by a reduced BDD, i.e. of the number of valuations ν s.t. $\nu \models f$.

Exercise 4 (Shared BDDs). When dealing with several boolean functions at once, with a fixed order on the variables, one can share the reduced BDDs for identical subfunctions. A *shared BDD* between m functions is a reduced BDD with m *root pointers* assigning a root node to each of the functions.

Let x_1, \dots, x_{2n} be the ordered set of variables. We want to compute the $n + 1$ bits $f_{n+1}f_n \cdots f_1$ of the sum of two n bits numbers $x_1x_3 \cdots x_{2n-1}$ and $x_2x_4 \cdots x_{2n}$. Represent the shared BDD for the functions f_3, \dots, f_1 , i.e. for $n = 2$.

Exercise 5 (An Upper Bound on the Size of BDDs). The size $B(f)$ of a reduced BDD for a function f is defined as the number of its nodes. Consider an arbitrary boolean function f on the ordered set $x_1 \cdots x_n$, and consider a variable x_k .

1. Show that we can bound the number of nodes labeled by $\{x_1, \dots, x_k\}$ by 2^k .
2. How many different subfunctions on the ordered set of variables $x_{k+1} \cdots x_n$ exist? Deduce another bound for the number of nodes labeled by $\{x_{k+1}, \dots, x_n\}$.

3. What global bound do you obtain for $k = n - \log_2(n - \log_2 n)$?

Exercise 6 (Finding the Optimal Order). There are in general $n!$ different orders for the variables $\{x_1, \dots, x_n\}$. One can nevertheless design an exponential time algorithm for finding the optimal order. Indeed, an optimal ordering on a subset X of variables does not depend on the order in which $\{x_1, \dots, x_n\} \setminus X$ has been accessed.

1. Given a subset X of $\{x_1, \dots, x_n\}$ and a variable x in X , how many nodes labeled by x does a BDD have if it first treats $\{x_1, \dots, x_n\} \setminus X$, then x , and last $X \setminus \{x\}$?
2. Reduce the optimal order problem to the search of a path of minimal weight in a weighted graph with subsets of $\{x_1, \dots, x_n\}$ as vertices.
3. Find the optimal order for the functions of Exercises 1.1 and 1.4.

Exercise 7 (Quasi Reduced BDDs). An ordered BDD for a boolean function f on $\{x_1, \dots, x_n\}$ is *complete* if all paths from the root to a sink are of length n . A BDD is *quasi reduced* if it is complete and no two nodes define the same subfunction.

1. Show that a quasi reduced BDD is unique up to isomorphism for an ordered set of variables $x_1 \cdots x_n$.
2. Let $Q(f)$ be the size of the quasi reduced BDD for the boolean function f on the ordered set of variables $x_1 \cdots x_n$. Show that $Q(f) \leq (n - 1)Q(f)$.

Exercise 8 (Minimal DFAs). A deterministic finite automaton \mathcal{A} recognizes a boolean function f on the ordered set of variables $x_1 \cdots x_n$ if $L(\mathcal{A}) = \{\nu \in \{0, 1\}^n \mid \nu \models f\}$, i.e. \mathcal{A} recognizes exactly the solutions of f .

What are the relations between the reduced BDD, the quasi reduced BDD, and the minimal DFA recognizing the same boolean function f on the ordered set of variables $x_1 \cdots x_n$?