

# Logics for AI: From Dreams to Formal Methods

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# Symbolic AI

- ▶ Symbolic AI: knowledge representation with symbols and rules manipulating symbols for reasoning.
- ▶ Traditionally, symbolic AI convenient for settings where the rules are clear cut, and you can easily obtain input and transform it into symbols.
- ▶ Logics for AI:
  - formal languages,
  - semantics,
  - proof procedures to reason about it.

## Subfields in IJCAI'2023 proceedings

- ▶ Agent-based and Multi-agent Systems
- ▶ AI Ethics, Trust, Fairness
- ▶ Computer Vision
- ▶ Constraint Satisfaction and Optimization
- ▶ Data Mining
- ▶ Game Theory and Economic Paradigms
- ▶ Humans and AI
- ▶ Knowledge Representation and Reasoning
- ▶ Machine Learning
- ▶ Multidisciplinary Topics and Applications
- ▶ Natural Language Processing
- ▶ Planning and Scheduling
- ▶ Robotics
- ▶ Search
- ▶ Uncertainty in AI
- ▶ etc.

# Subfields in KR'23 proceedings

- ▶ Argumentation
- ▶ Automated reasoning
- ▶ Belief merging / revision
- ▶ Conditionals
- ▶ Description logics
- ▶ Epistemic logic
- ▶ Knowledge representation and machine learning
- ▶ Multi-agent systems
- ▶ Strategic reasoning
- ▶ Systems and robotics
- ▶ Temporal reasoning

## Reasoning on Ontologies

- ▶ Ontology: formal specification of some domain with concepts, objects, relationships between concepts, objects, etc.
- ▶ Backbone of ontologies includes:
  - taxonomy (classification of objects),
  - axioms (to constrain the models of the defined terms).
- ▶ Well-known ontologies:
  - Medical ontology SNOMED-CT.
  - NCI Thesaurus (National Cancer Institute, USA).
  - Gene ontology (world largest source of information on the functions of genes).

*(classification of medical terms: disease  
drug)*

- ▶ Free ontology editor Protégé

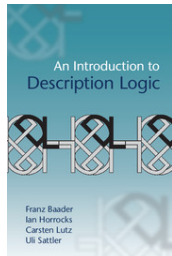
<http://protege.stanford.edu/>

# Challenges with Ontologies

- ▶ How to define ontologies and to reason on it?
- ▶ How to repair faulty ontologies ?
- ▶ How to add new concepts or axioms without affecting the old inferences?
- ▶ More generally, how to extract from the ontologies more knowledge than what is explicitly specified?
  - inferences about individuals,
  - concept subsumptions, non-redundancy,
  - concept hierarchy, consistency of concepts,
  - etc.

# Why Description Logics?

- ▶ Formal languages for concepts, relations and instances.
- ▶ DLs have all one needs to formalise ontologies.
- ▶ Computational properties.
  - Acceptable trade-off between expressivity and complexity.
  - Decidability and often tractability.
  - Implementation in tools of the main reasoning tasks.
- ▶ A remarkable suite of languages and tools.  
See e.g.,
  - OWL: Web Ontology Language.
  - Protégé: ontology editor.
  - FaCT++: DL reasoner supporting OWL DL.



## Description Logic $\mathcal{ALC}$ in a Nutshell

- ▶ Language of complex concepts.

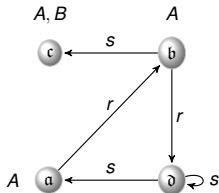
$$C ::= \top \mid \perp \mid A \mid \neg C \mid C \sqcap C \mid C \sqcup C \mid \exists r.C \mid \forall r.C,$$

with concept names  $A$  and role names  $r$ .

- ▶ Interpretation  $\mathcal{I} \stackrel{\text{def}}{=} (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$

- $\Delta^{\mathcal{I}}$ : non-empty set (the **domain**).
- $\cdot^{\mathcal{I}}$ : **interpretation function** such that

$$A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \quad r^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$$



- ▶  $C^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$  defined inductively providing the semantics to complex concepts.



# Concept Inclusions and Decision Problems

- ▶ General concept inclusions  $C \sqsubseteq D$  (GCIs).

E.g.,  $\text{Employee} \sqsubseteq \exists \text{WorksFor. T}$

$$\mathcal{I} \models C \sqsubseteq D \stackrel{\text{def}}{\iff} C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$$

- ▶ Terminological Box (TBox)  $\mathcal{T}$ : finite set of GCIs.
- ▶ Interpretation  $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$ , TBox  $\mathcal{T}$ .

$$\mathcal{I} \models \mathcal{T} \stackrel{\text{def}}{\iff} \text{for all } C \sqsubseteq D \in \mathcal{T}, \mathcal{I} \models C \sqsubseteq D$$

- ▶ **Concept satisfiability problem w.r.t. general TBoxes:**

**Input:** A concept  $C_0$  and a TBox  $\mathcal{T}$ .

**Question:** Is there an interpretation  $\mathcal{I}$  such that  $\mathcal{I} \models \mathcal{T}$   
and  $C_0^{\mathcal{I}} \neq \emptyset$ ?

- ▶ This problem is EXPTIME-complete.

# Description Logics with Concrete Domains

- ▶ Need to express concrete properties about data in ontologies (e.g. age, duration, name, size, etc.)
- ▶ Examples of concrete domains:  
 $(\mathbb{N}, <, +1)$ ,  $(\mathbb{Q}, <, =)$ ,  $(\mathbb{N}, <, =)$ ,  $(\{0, 1\}^*, <_{\text{pre}}, <_{\text{suf}})$ .
- ▶ General scheme for integrating concrete domains in DLs.  
[Baader & Hanschke, IJCAI'91]
  - declarative semantics close to the usual semantics for DLs,
  - generic extensions of DLs with various concrete domains,
  - tableaux-based algorithms combined with theory reasoning.

# Methods for Handling Concrete Domains

- ▶ Tableaux-based **calculi** for “ $\omega$ -admissible” domains.  
[Lutz & Miličić, JAR 2007]
- ▶ Translation into decidable **logics** [Carapelle & Turhan, ECAI'16]
  - Decidability of concept satisfiability problem w.r.t. general TBoxes for  $\mathcal{ALC}(\mathbb{N}, <, =, (=_n)_{n \in \mathbb{N}})$ .
- ▶ Translation into **automata**-based problems.
  - Concept satisfiability problem w.r.t. general TBoxes for  $\mathcal{ALC}(\mathbb{N}, <, =, (=_n)_{n \in \mathbb{N}})$  in EXPTIME.  
[Labai & Ortiz & Šimkus, KR'20]
  - Concept satisfiability problem w.r.t. general TBoxes for  $\mathcal{ALC}(\{0, 1\}^*, <_{\text{pre}})$  in EXPTIME. [Demri & Quaas, JELIA'23]

*(constraint automata for data words or data tree)*

# Temporal Logics with Concrete Domains

- ▶ Concrete domains in TCS:
  - Satisfiability Modulo Theory (SMT) solvers.  
String theories, arithmetical theories, array theories, etc.
  - Verification of database-driven systems.
  - Temporal logics with arithmetical constraints.
- ▶ “Infinitely often  $x$  is a prefix of the next value for  $y$ ”:  
 $GF(x <_{\text{pre}} Xy)$ .
- ▶ Satisfiability problem for  $\text{CTL}^*(\mathbb{Z}, <, =, (=_n)_{n \in \mathbb{Z}})$  is decidable in  $2\text{EXPTIME}$ .

[Carapelle et al, JCSS 2016; Demri & Quaas, CONCUR'23]

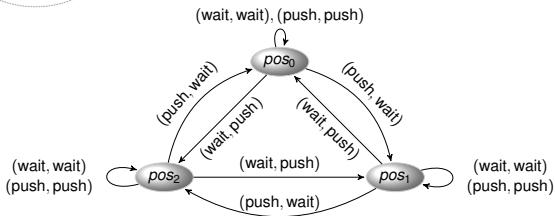
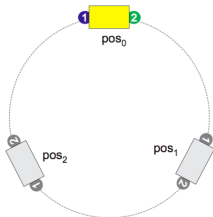
(  $\text{CTL}^*$ : well-known logic related to model-checking)

## Another Success Story: Logics of Strategic Ability

- ▶ To express that a coalition of agents has a collective strategy to achieve some goal and to reason on it.
- ▶ A strategy is a conditional plan intended to work whatever the other agents do.
- ▶ Well-known specimens.
  - Coalition Logic CL. (one-step strategies)
  - Alternating-time temporal logic ATL. (generalisation of temporal logics)
  - Strategy Logic SL. (explicit quantification over strategies)  
*(with a huge amount of variants)*

# Multi-Agents Systems

- Multi-agent systems are transition systems in which transitions are fired when simultaneous actions are performed by different agents.



## ATL-like logics

- ▶  $\langle\langle A \rangle\rangle \Phi$ : coalition  $A$  has a collective strategy to enforce the temporal property  $\Phi$ .
- ▶ A collective strategy is a tuple of individual strategies.

$$\varphi ::= p \mid \neg \varphi \mid \varphi \wedge \varphi \mid \langle\langle A \rangle\rangle X\varphi \mid \langle\langle A \rangle\rangle G\varphi \mid \langle\langle A \rangle\rangle \varphi U \varphi$$
$$p \in \text{PROP} \quad A \subseteq \text{Agt}$$

- ▶  $\mathfrak{M}, s \models \langle\langle A \rangle\rangle G\varphi \stackrel{\text{def}}{\iff} \exists \text{ strategy } \sigma \text{ such that}$ 
  - $\forall \text{ computations } \lambda = s_0 \rightarrow s_1 \dots \text{ from } s \text{ respecting } \sigma,$
  - $\forall \text{ positions } i, \text{ we have } \mathfrak{M}, s_i \models \varphi.$
- ▶ Tractable model-checking problems and automata-based satisfiability-checking decision procedures.

## More Ingredients

▶ Resource-aware logics:

- actions have costs/weights,
- formulae may specify constraints about such (cumulative) costs/weights.

E.g. [Belardinelli & Demri, AI 2021; Bulling & Goranko, AAMAS 2022]

*(relationships with energy game)*

▶ Strategic reasoning with knowledge operators.

See e.g. [Agotnes, Synthese 2006]

▶ Restriction on agents' knowledge.

- Strategy logics with imperfect information.
- Undecidability can be obtained easily.
- Fragments including those with hierarchies of knowledge leads to less expensive reasoning tasks.

See e.g. [Berthon et al, TOCL 2021]



## In My Reading List

- ▶ Well-identified potential interactions between machine learning and symbolic AI.
  - Machine learning can be used to solve logical problems and to accelerate verification/automated techniques.
  - Logical methods can be used to complement learning algorithms to improve the precision and explainability.

*(hybridization?)*

- ▶ Unique characterisability and learnability of temporal instance queries

[Fortin et al., KR'22]

*(Example set  $(E^+, E^-)$  to characterise temporal formulae)*

- ▶ An SMT-Based approach for verifying binarized neural networks

[Amir et al., TACAS'21]

*(SMT-based approach for formal verification of binarized neural networks)*

# Logics for AI: the Great Return ?

- ▶ Beyond knowledge representation and reasoning for description logics, strategy logics, dynamic epistemic logics, etc. ? *(new applications?)*
- ▶ Arithmetical theories for the verification of neural networks. *(new arithmetical theories)*
- ▶ Learning logical formulae and unique characterisations. *(new logical problems?)*

## **A Selection of Bibliographical References**

# Logical Formalisms with Concrete Domains

- ▶ “Description Logics with Concrete Domains—A Survey”

[Lutz, AiML'02]

*(a classical paper on de  
domains)*

- ▶ “Concrete domains in logics: a survey”

[Demri & Quaas, SIGLOG News 2021]

*(a brief survey)*

- ▶ “Using Model Theory to Find Decidable and Tractable Description Logics with Concrete Domains”

[Baader & Rydval, JAR 2022]

*(model theory for DL)*

# Strategic Reasoning and Resources

- ▶ “Alternating-time temporal logic”

[Alur & Henzinger & Kupferman, JACM 2022]

*(the classical paper about ATL)*

- ▶ “Combining quantitative and qualitative reasoning in concurrent multi-player games”

[Bulling & Goranko, AAMAS 2022]

*(how to mix quantitative and qualitative ob)*

- ▶ “Strategic reasoning with a bounded number of resources:  
The quest for tractability”

[Belardinelli & Demri, AI 2021]

*(complexity analy*

**ATL)**

# Learning and Modal/Temporal Logics

- ▶ “Scalable Anytime Algorithms for Learning Fragments of Linear Temporal Logic” [Raha et al., JOSS 2024]

*(how to learn temporal formulae)*

- ▶ Unique characterisability and learnability of temporal instance queries [Fortin et al., KR'22]

*(Example set  $(E^+, E^-)$  to characterise temporal formulae)*

- ▶ Logic of “Black Box” classifier systems [Liu & Lorini, WoLLIC'22]

*(product modal logic for multi-classifier model)*

# Logic, Verification and Neural Networks

- ▶ “Neural Network Verification with Proof Production”

[Isac et al., FMCAD'22]

*(how to add proof production capabilities)*

- ▶ “An SMT-Based approach for verifying binarized neural networks”

[Amir et al., TACAS'21]

*(SMT-based approach for formal verification of binarized neural networks)*

- ▶ “Simplifying neural networks using formal verification”

[Gokulanathan et al., NFM'20]

*(how to remove component)*