

Probabilistic Aspects of Computer Science: Exercise 5

Exercise 1 Given an alphabet Σ , a PFA \mathcal{A} on Σ , and two numbers $x, y \in (0, 1)$, show that there is a \mathcal{B} such that $\mathcal{L}_{>x}(\mathcal{A}) = \mathcal{L}_{>y}(\mathcal{B})$ and $\mathcal{L}_{\geq x}(\mathcal{A}) = \mathcal{L}_{\geq y}(\mathcal{B})$.

Exercise 2 Given an alphabet Σ and a PFA \mathcal{A} on Σ :

1. For any regular language $L \subset \Sigma^*$, there is a PFA \mathcal{A} such that $L = \mathcal{L}_{=1}(\mathcal{A})$.
2. The problem of checking emptiness of $\mathcal{L}_{=1}(\mathcal{A})$ is PSPACE-complete.
3. The problem of checking universality of $\mathcal{L}_{=1}(\mathcal{A})$ is decidable in polynomial time.

Exercise 3 Let $\Sigma = \{\mathbf{0}, \mathbf{1}\}$. Consider the PFA $\mathcal{A} = (Q, q_s, Q_f, \delta)$ where the set $Q = (q_s, q_a, q_r)$, $Q_f = \{q_a\}$ and δ is defined as follows.

- $\delta(q_s, \mathbf{0}, q_s) = \delta(q_s, \mathbf{0}, q_r) = \frac{1}{2}$.
- $\delta(q_s, \mathbf{1}, q_s) = \delta(q_s, \mathbf{1}, q_a) = \frac{1}{2}$.
- $\delta(q_a, \mathbf{0}, q_a) = \delta(q_a, \mathbf{1}, q_a) = 1$.
- $\delta(q_r, \mathbf{0}, q_r) = \delta(q_r, \mathbf{1}, q_r) = 1$.

1. Show that for any word $u = a_1 \cdots a_n \in \Sigma^*$, the probability of \mathcal{A} accepting u is

$$\sum_{i=1}^n \frac{\text{num}(a_i)}{2^i}$$

where $\text{num}(a_i)$ is the number 0 if a_i is $\mathbf{0}$ and is number 1 one when a_i is $\mathbf{1}$.

2. Compute $\mathcal{L}_{>\frac{1}{4}}(\mathcal{A})$. Is it regular?

Exercise 4 Given any PFA \mathcal{A} , show that there is another automata \mathcal{B} such that:

1. $\mathcal{L}_{>\frac{1}{2}}(\mathcal{B}) = \mathcal{L}_{>\frac{1}{\sqrt{2}}}(\mathcal{A})$.
2. In general for any $x \in [0, 1]$, $\mathcal{L}_{>x^2}(\mathcal{B}) = \mathcal{L}_{>x}(\mathcal{A})$.