Schnorr's Protocol

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Summary: The Schnorr protocol is described by R. Cramer, I. Damgård and B. Schoenmakers in [CDS94].

Protocol specification (in common syntax)

A, B: principal Na, Nb: fresh number Sa : private key Pa = exp(g,Sa) : public key A chooses Na and computes a = exp(g, Na)1. Α -> B : а B chooses Nb 2. B -> A : Nb A computes r = Na + Nb \times Sa 3. A -> B : r B checks that $\exp(g, r) = a \times \exp(Pa, Nb)$

Description of the protocol rules

A zero-knowledge protocol is designed for convincing the verifier of the validity of a given statement, without releasing any knowledge beyond the validity of the statement. This concept was introduced in [GMR85]. An overview can be found in [Gol01]. We present the Schnorr protocol which is described by R. Cramer, I. Damgård and B. Schoenmakers in [CDS94] and which uses this method.

The +, \times and exp symbols denote respectively addition, multiplication and modular exponentiation.

Details of the computation done by B at the last step of the protocol:

 $a \times exp(Pa,Nb)$ $= exp(g,Na) \times exp(exp(g,Sa),Nb)$ $= exp(g,Na) \times exp(g,Sa \times Nb)$ $= exp(g,Na + Sa \times Nb)$ = exp(g, r)

Requirements

 ${\tt A}$ wants to prove his identity to ${\tt B}$ by showing him that he knows ${\tt Sa}$ without revealing it.

References

[CDS94]

Citations

- [CDS94] R. Cramer, I. Damgård, and B. Schoenmakers. Proofs of partial knowledge and simplified design of witness hiding protocols. In Proc. 14th Annual International Cryptology Conference (CRYPTO'94), volume 963 of LNCS, pages 174–187, Santa Barbara (California, USA), 1994. Springer-Verlag.
- [GMR85] S. Goldwasser, S. Micali, and C. Rackoff. The knowledge complexity of interactive proof-systems. In Proc. 17th annual ACM Symposium on Theory of Computing, pages 291–304. ACM Press, 1985.
- [Gol01] O. Goldreich. Foundations of Cryptography. Cambridge University Press, 2001.