

Diffie Helman

Author(s): W. Diffie and M. Helman 1978

Last modified November 11, 2002

Summary: The Diffie Helman key exchange algorithm.

Protocol specification (in common syntax)

```

A, B :          principal
P, G, Xa, Xb :  number
one :          -> number
kap :          number, number, number -> number

1.   A -> B :    P, G
2.   A -> B :    kap(P, G, Xa)
3.   B -> A :    kap(P, G, Xb)
4.   A -> B :    {one()}kap(P, kap(P, G, Xb), Xa)

```

Description of the protocol rules

The function `kap` must satisfy:

$$\text{kap}(P, \text{kap}(P, G, Y), X) = \text{kap}(P, \text{kap}(P, G, X), Y)$$

It is implemented by: $\text{kap}(P, X, Y) = \exp(X, Y) \bmod P$.

It the protocol, `P` is choosen to be a prime number `P` and `G < P`.

The fresh symmetric key exchanged is $\text{kap}(P, \text{kap}(P, G, Xb), Xa) = \text{kap}(P, \text{kap}(P, G, Xa), Xb)$.

Requirements

The protocol must guaranty the secrecy of the fresh key.

The protocol must guaranty the authenticity of the participants.

References

[DH76]

Claimed proofs

[Bla01]

Claimed attacks

The authenticity is not guaranteed by the protocol.

- | | | | | | | |
|----|------|----|------|---|----------------------------------|----|
| 1. | I(A) | -> | B | : | P, G | |
| 2. | I(A) | -> | B | : | kap(P, G, Xi) | |
| 3. | B | -> | I(A) | : | kap(P, G, Xb) | or |
| 4. | I(A) | -> | B | : | {one()}kap(P, kap(P, G, Xb), Xi) | |
| | | | | | | |
| 1. | A | -> | I(B) | : | P, G | |
| 2. | A | -> | I(B) | : | kap(P, G, Xa) | |
| 3. | I(B) | -> | A | : | kap(P, G, Xi) | |
| 4. | A | -> | I(B) | : | {one()}kap(P, kap(P, G, Xi), Xa) | |

Citations

- [Bla01] Bruno Blanchet. An efficient cryptographic protocol verifier based on prolog rules. In IEEE, editor, *14th IEEE Computer Security Foundations Workshop (CSFW-14)*, june 2001.
- [DH76] W. Diffie and M. Helman. New directions in cryptography. *IEEE Transactions on Information Society*, 22(6):644–654, november 1976.